

## Linear Least Squares

Given a linear system  $\mathbf{Ax} - \mathbf{b} = \mathbf{e}$ ,

$$\mathbf{a}_1 \bullet \mathbf{x} - b_1 = e_1$$

$\vdots$

$$\mathbf{a}_i \bullet \mathbf{x} - b_i = e_i$$

$\vdots$

$$\mathbf{a}_m \bullet \mathbf{x} - b_m = e_m$$

We want to minimize the sum of squares of the errors

$$\min_{\mathbf{x}} \sum_i e_i^2 = \mathbf{e}^T \mathbf{e} = (\mathbf{Ax} - \mathbf{b})^T (\mathbf{Ax} - \mathbf{b})$$

Sometimes write this as  $\mathbf{Ax} \cong \mathbf{b}$

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## Linear Least Squares

- Many methods for  $\mathbf{Ax} \approx \mathbf{b}$
- One simple one is to compute

$$\mathbf{Ax} \approx \mathbf{b}$$

$$\mathbf{A}^T \mathbf{Ax} \approx \mathbf{A}^T \mathbf{b}$$

$$\mathbf{x} \approx (\mathbf{A}^T \mathbf{A})^{-1} \mathbf{A}^T \mathbf{b}$$

- Better methods based on orthogonal transformations exist
- These methods are available in standard math libraries
- A short review follows

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## Orthogonal Transformations

The key property is:

$$\mathbf{Q}^{-1} = \mathbf{Q}^T$$

Some implications of this are as follows

$$\text{if } \mathbf{Q} = [\mathbf{q}_1 \quad \mathbf{q}_2 \quad \cdots \quad \mathbf{q}_n]$$

$$\text{then } \mathbf{q}_i \bullet \mathbf{q}_j = \begin{cases} 1, & \text{if } i=j \\ 0, & \text{if } i \neq j \end{cases}$$

$$\|\mathbf{Q}\mathbf{x}\| = \sqrt{(\mathbf{Q}\mathbf{x})^T (\mathbf{Q}\mathbf{x})}$$

$$= \sqrt{\mathbf{x}^T \mathbf{Q}^T \mathbf{Q} \mathbf{x}} = \sqrt{\mathbf{x}^T \mathbf{x}}$$

$$= \|\mathbf{x}\|$$

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## General Approach

The discussion below generally follows the development in

D. Lawson and R. Hanson, *Solving Least Squares Problems*,  
Prentice-Hall, 1974

However, similar discussions may be found in many textbooks.

Given the problem

$$\min \|\mathbf{Ax} - \mathbf{b}\|$$

Observe that for any orthogonal matrix  $\mathbf{Q}$

$$\|\mathbf{Ax} - \mathbf{b}\| = \|\mathbf{Q}(\mathbf{Ax} - \mathbf{b})\| = \|\mathbf{QAx} - \mathbf{Qb}\|$$

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## Theorem (from Lawson & Hanson pp 5-6)

Suppose  $\mathbf{A}$  is an  $m \times n$  matrix with rank  $k$  and

$$\mathbf{A} = \mathbf{H}\mathbf{R}\mathbf{K}^T$$

where

$\mathbf{H} = m \times m$  orthogonal matrix

$\mathbf{K} = n \times n$  orthogonal matrix

$$\mathbf{R} = \begin{bmatrix} \mathbf{R}_{11} & \mathbf{0} \\ \mathbf{0} & \mathbf{0} \end{bmatrix} \text{ with } \text{rank}(\mathbf{R}_{11}) = k$$

This is called an orthogonal decomposition of  $\mathbf{A}$

Define

$$\mathbf{g} = \mathbf{H}^T \mathbf{b} = \begin{bmatrix} \mathbf{g}_1 \\ \mathbf{g}_2 \end{bmatrix} \begin{matrix} \} k \\ \} n-k \end{matrix} \quad \mathbf{y} = \mathbf{K}^T \mathbf{x} = \begin{bmatrix} \mathbf{y}_1 \\ \mathbf{y}_2 \end{bmatrix} \begin{matrix} \} k \\ \} n-k \end{matrix}$$

and define  $\tilde{\mathbf{y}}_1$  to be the unique solution of

$$\mathbf{R}_{11} \mathbf{y}_1 = \mathbf{g}_1$$

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## Theorem (from Lawson & Hanson pp 5-6)

Then ...

- 1) All solutions to the problem of minimizing  $\|\mathbf{Ax} - \mathbf{b}\|$  are of the form

$$\hat{\mathbf{x}} = \mathbf{K} \begin{bmatrix} \tilde{\mathbf{y}}_1 \\ \mathbf{y}_2 \end{bmatrix} \text{ where } \mathbf{y}_2 \text{ is arbitrary}$$

- 2) Any such  $\hat{\mathbf{x}}$  produces the same residual vector  $\mathbf{r}$  satisfying

$$\mathbf{r} = \mathbf{b} - \mathbf{A}\hat{\mathbf{x}} = \mathbf{H} \begin{bmatrix} \mathbf{0} \\ \mathbf{g}_2 \end{bmatrix}$$

- 3) The norm of  $\mathbf{r}$  satisfies

$$\|\mathbf{r}\| = \|\mathbf{b} - \mathbf{A}\hat{\mathbf{x}}\| = \|\mathbf{g}_2\|$$

- 4) The unique solution of minimum length is

$$\tilde{\mathbf{x}} = \mathbf{K} \begin{bmatrix} \tilde{\mathbf{y}}_1 \\ \mathbf{0} \end{bmatrix}$$

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## Householder Decomposition

One method uses repeated Householder transformations to produce an upper triangular matrix  $\mathbf{R}$ .

$$\mathbf{H}^T \mathbf{A} \mathbf{K} = \mathbf{R} = \begin{bmatrix} \mathbf{R}_{11} & \mathbf{0} \\ \mathbf{0} & \mathbf{0} \end{bmatrix} = \begin{bmatrix} r_{11} & r_{12} & \cdots & r_{1k} & 0 & \cdots & 0 \\ 0 & r_{22} & \ddots & \vdots & \vdots & \ddots & \vdots \\ \vdots & & \ddots & r_{k-1,k} & \vdots & & \vdots \\ 0 & \cdots & 0 & r_{kk} & 0 & \cdots & 0 \\ 0 & \cdots & 0 & 0 & 0 & \cdots & 0 \\ \vdots & \ddots & \vdots & \vdots & \vdots & \ddots & \vdots \\ 0 & \cdots & 0 & 0 & 0 & \cdots & 0 \end{bmatrix}$$

where  $\mathbf{H}^T = \mathbf{H}_{k-1}^T \cdots \mathbf{H}_2^T \mathbf{H}_1^T$  is a product of Householder transformations and  $\mathbf{K} = \mathbf{K}_1 \mathbf{K}_2 \cdots \mathbf{K}_p$  is a series of permutations, if needed, to avoid division by 0. Then, we solve the problem  $\mathbf{A} \mathbf{x} \approx \mathbf{b}$  by solving  $\mathbf{R}_{11} \tilde{\mathbf{y}}_1 = \mathbf{g}_1$  and

forming  $\tilde{\mathbf{x}} = \mathbf{K} \begin{bmatrix} \tilde{\mathbf{y}}_1 \\ \mathbf{0} \end{bmatrix}$  as outlined before.

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## Singular Value Decomposition

- Developed by Golub, et al in late 1960's
- Commonly available in mathematical libraries
- E.g.,
  - MATLAB
  - IMSL
  - Numerical Recipes (Wm. Press, et. al., Cambridge Press)
  - CISST ERC Math Library

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## Singular Value Decomposition

Given an arbitrary  $m$  by  $n$  matrix  $\mathbf{A}$ , there exist orthogonal matrices  $\mathbf{U}$ ,  $\mathbf{V}$  and a diagonal matrix  $\mathbf{S}$  that:

$$\mathbf{A}_{m \times n} = \mathbf{U}_{m \times m} \begin{bmatrix} \mathbf{S}_{n \times n} \\ \mathbf{0}_{(m-n) \times n} \end{bmatrix} \mathbf{V}_{n \times n}^T \quad \text{for } m \geq n$$

or

$$\mathbf{A}_{m \times n} = \mathbf{U}_{m \times m} \begin{bmatrix} \mathbf{S}_{n \times n} & \mathbf{0}_{(m \times (n-m))} \end{bmatrix} \mathbf{V}_{n \times n}^T \quad \text{for } m \leq n$$

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## SVD Least Squares

$$\mathbf{A}_{m \times n} \mathbf{x} \approx \mathbf{b}$$

$$\mathbf{U}_{m \times m} \begin{bmatrix} \mathbf{S}_{n \times n} \\ \mathbf{0}_{(m-n) \times n} \end{bmatrix} \mathbf{V}_{n \times n}^T \mathbf{x} = \mathbf{b}$$

$$\begin{bmatrix} \mathbf{S}_{n \times n} \\ \mathbf{0}_{(m-n) \times n} \end{bmatrix} \mathbf{y} = \mathbf{U}_{m \times m}^T \mathbf{b} \quad \text{where } \mathbf{y} = \mathbf{V}^T \mathbf{x}$$

Solve this for  $\mathbf{y}$  (trivial, since  $\mathbf{S}$  is diagonal), then compute

$$\mathbf{V} \mathbf{y} = \mathbf{V} \mathbf{V}^T \mathbf{x} = \mathbf{x}$$

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## Least squares adjustment

Given a vector function  $\vec{\mathbf{G}}(\vec{\mathbf{q}}; \vec{\mathbf{u}})$  of parameters  $\vec{\mathbf{q}}$  and experimental variables  $\vec{\mathbf{u}}$ , together with a set of observations

$$\vec{\mathbf{v}}_k = \vec{\mathbf{G}}(\vec{\mathbf{q}}; \vec{\mathbf{u}}_k)$$

and an initial guess  $\vec{\mathbf{q}}_0$  of the values of  $\vec{\mathbf{q}}$ , we wish to find a better estimate of  $\vec{\mathbf{q}}$ .

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## Least Squares Adjustment

Step 0  $j \leftarrow 0$ ;

Step 1 Compute  $\vec{\mathbf{e}}_k \leftarrow \vec{\mathbf{v}}_k - \vec{\mathbf{G}}(\vec{\mathbf{q}}_j; \vec{\mathbf{u}}_k)$ ;  $\vec{\mathbf{E}}_j \leftarrow [\vec{\mathbf{e}}_1, \dots, \vec{\mathbf{e}}_N]^T$

Step 2 If  $\|\vec{\mathbf{E}}_j\|$  is small or some other convergence criterion is met, then stop. Otherwise go on to Step 3.

Step 3 Solve the least squares problem

$$\begin{bmatrix} \vdots \\ \mathbf{J}_G(\vec{\mathbf{q}}_j, \vec{\mathbf{u}}_k) \\ \vdots \end{bmatrix} \bullet \Delta \vec{\mathbf{q}} \approx \begin{bmatrix} \vdots \\ \vdots \\ \vdots \end{bmatrix}$$

for  $\Delta \vec{\mathbf{q}}$ .

Step 4 Set  $\vec{\mathbf{q}}_{j+1} \leftarrow \vec{\mathbf{q}}_j + \Delta \vec{\mathbf{q}}$ ;  $j \leftarrow j + 1$ ; Go back to Step 1.

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